

Event-B Development of a Smart Ballot Box

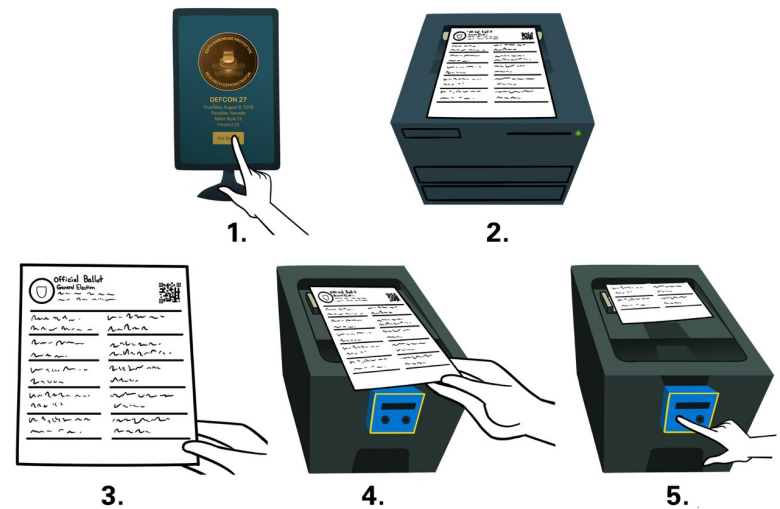
Dana Dghaym

Outline

- **Part 1** : Event-B development of a smart ballot box
 - Case study: Smart Ballot Box
 - Aims & Motivation
 - Event-B System Model of a Smart Ballot Box (SBB)
- **Part 2** : Event-B to SPARK-Ada
 - Introduction to SPARK
 - High-Level Transformation Patterns : SBB Example
 - Transformation Issues
- Conclusions and Future Work

Case Study: Smart Ballot Box

- Key Function of SBB:
 - Ensures only valid ballot papers are cast in ballot boxes for later tabulation.
- Why the SBB?
 - Security Properties: Confidentiality, integrity and availability.
 - Relatively Small Case Study
 - Can be used as Demonstrator



Galois and Free & Fair. The BESSPIN Voting System (2019).

Motivation

- Application of refinement-based formal modelling in building a **Correct-by-Construction** secure system.
- Refinement of the **security properties** of a system.
- **Overall Aim** of case study:
 - Show how the Smart Ballot Box can be correctly implemented on **capability hardware** according to the system-level security specification.
- Develop a tool-supported approach to translate an Event-B model to verified code.

SBB System Model: Refinement Strategy

0. Abstract level: Model an **ideal voting system**.
1. Model possible **attackers'** behavior by distinguishing between different types of ballot papers.
2. Introduce **time** and invalidate ballots with expired timestamps.
 - Time can be the subject of more attacks.
3. Data refine the voter information by **encrypting** ballots.
4. Ensure the legitimacy of ballots through the Message Authentication Code (**MAC**).

SBB Model: Abstract Level

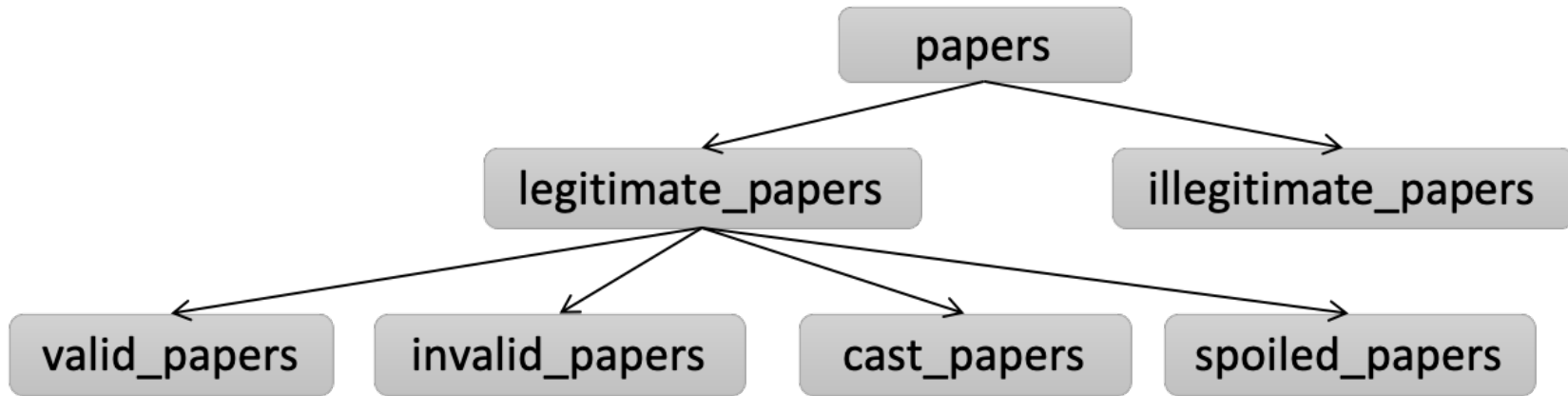
- Events: create_ballot, cast ballot, invalidate_ballot
- Model an ideal voting system
 - Each voter can have at most one legitimate ballot

ballots \in VOTER \rightarrow VOTE

- The cast ballots must be legitimate

cast \subseteq ballots

First Refinement: Ballot types



- Possible attacks
 - Attacker create ballot/duplicate valid ballot ..
- Model the main **security** properties of SBB
 1. Accept all valid ballots
 2. Reject invalid ballots

First Refinement: Availability Property

- Availability property: Ensure valid ballots are not blocked from being cast.
- Availability property is captured by the guard of the relevant events.
 - Specify *cast_paper* as **rigid event** with *paper* as **rigid parameter** .

***Rigid:** the guard cannot be strengthened

```
event [cast_paper] refines cast_ballot
any [paper] where
    @valid-paper: paper ∈ valid_papers
then
    // actions for casting a ballot
end
```


Second Refinement: Time & Availability

- This theorem ensures that a ballot paper is considered **valid** only if:
 - Paper time has **not expired**
 - Voter has **not cast** their vote before
 - The paper is **not spoiled**
 - Issued by a **legitimate** source

```
theorem @accept-valid-paper:  
  ∀ paper · paper ∈ valid_papers ⇒  
    // paper not already expired  
    paper_time(paper) ≥ current_time -  
    expiry_duration  
    // copy not already cast  
    ∧ paper_voter(paper) ∉  
    paper_voter[cast_papers]  
    // copy not already spoiled  
    ∧ (∀ sp · sp ∈ spoiled_papers ⇒  
      paper_voter(paper) ≠ paper_voter(sp)  
      ∧ paper_vote(paper) ≠ paper_vote(sp)  
      ∧ paper_time(paper) ≠ paper_time(sp)  
    )  
    // paper is not illegitimate  
    ∧ paper ∉ illegitimate_papers
```

Third Refinement: Ballot Encryption

- Introduce encryption to prevent SBB from accessing the voter's information.
 - Apply **data refinement** to replace *paper_voter* and *paper_vote* with encrypted ballot

```

theorem @accept-valid-paper:
   $\forall$  paper  $\cdot$  paper  $\in$  valid_papers  $\Rightarrow$ 
    paper_time(paper)  $\geq$  current_time -
    expiry_duration
  // copy not already cast
   $\wedge$  paper_encrypted_ballot(paper)  $\notin$ 
  paper_encrypted_ballot[cast_papers]
  // copy not already spoiled
   $\wedge$  ( $\forall$  sp  $\cdot$  sp  $\in$  spoiled_papers  $\Rightarrow$ 
    paper_encrypted_ballot(paper)  $\neq$ 
    paper_encrypted_ballot(sp)  $\vee$ 
    paper_time(paper)  $\neq$  paper_time(sp)
  )
   $\wedge$  paper  $\notin$  illegitimate_papers

```

Fourth Refinement: Ballot Authentication

- Introduce **MAC** to check the legitimacy of the source issuing the ballot.
 - We assume the attacker does not know the secret key; therefore, it is crucial to ensure the secrecy of this key.

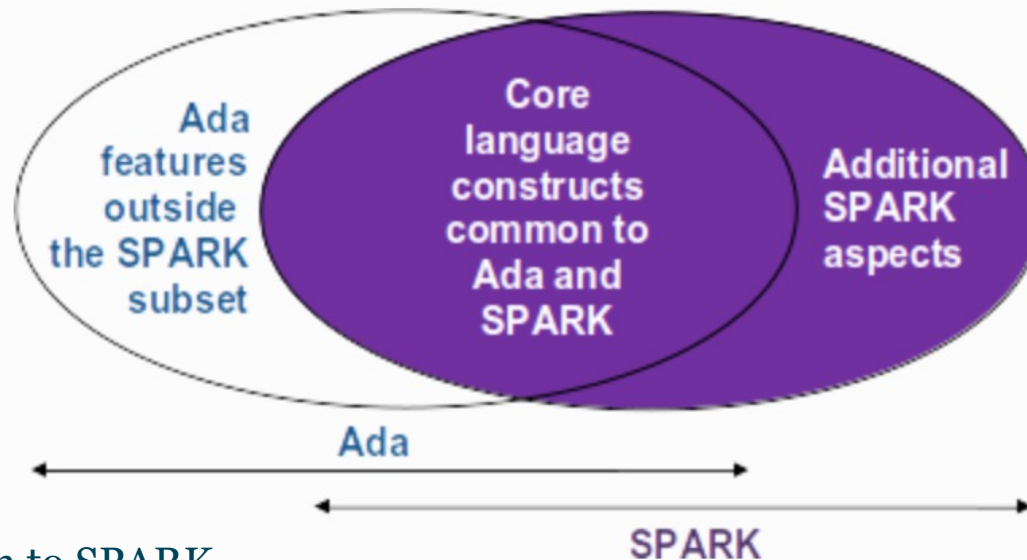
```
@mac-legitimate_papers:  $\forall$ paper · paper  $\in$  legitimate_papers  $\Rightarrow$   
paper_mac(paper) = MACAlgorithm(  
paper_time(paper)  $\mapsto$  paper_encrypted_ballot(paper)  $\mapsto$  MACKey  
)
```

- All proofs are automatically proved with the help of SMT-Solver plugin

Part 2: Translating Event-B to SPARK- Ada

Introduction

- What is **SPARK**?
 - A programming language based on a **subset** of the Ada language,
 - Targeted at functional **specification** and static **verification**.
 - A set of development and verification tools for that language.



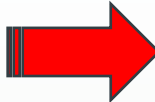
From Event-B to SPARK

```

machine m
sees C
variables
...
invariants
...
events
  event
INITIALISATION
  then
    @act1: ...
  end
event evt
  any
    parameters
  where
    @grd1: ..
  then
    @act1: ...
  end
end
  
```

```

context C
sets
...
constants
...
records
  record T
    A : Integer
    B : Integer
  ...
axioms
...
end
  
```



```

package P
with SPARK_Mode => On
is
Some_Global : G;
type T is record
  ...
  A : Integer;
  B : Integer
end record;

function F(X :)
return ..
...
;
procedure Proc (X : in T)
with
  Global => (Input => (..,
    ... )),
  Pre => .. ,
  Post => .. ;
end P;
  
```

```

package body P
with SPARK_Mode => On
is
procedure Proc (X : in T) is
begin
  ... ;
end Proc;
end P;
  
```

Refinement towards Implementation

- Sets \rightarrow arrays
 - Data Refine a set to **Total function** from Integer range to the set type
 - Can introduce a counter variable to track the size of array
- Event-B records are more general than SPARK (Event-B records supports optional and relational fields)
 - Use only total functions
 - Define a **special null record** element to reflect Event-B optional possibility and it can be used for initialisations

High level Event-B Transformation

- **Event-B Models Translation**
 - Each context → specification package using all extended context packages
 - Last Refined Machine → specification and body packages using all context and extended contexts packages
- **Machine Elements Translation**
 - Variables → Global variables, initialised according to the INITIALISATION event actions
 - Event / INITIALISATION → Procedures
 - Event Guards → Pre-conditions
 - Event Actions → Post-conditions
 - Event Parameters → Procedure Parameters (Output, input, in out depends on guards and actions)

Smart Ballot Example

```

event cast_paper
refines cast_paper
any
paper
where
@grd1: paper ∈ BARCODE
@grd2: cast_count ∈ 0 .. max_votes - 1
      ....
then
@act1: cast_arr(cast_count) := paper
@act2: cast_count := cast_count + 1
end

```

```

procedure cast(paper : in barcode) with
Global => (Proof_In => (spoiled_arr,
curr_time, spoil_count),
In_Out => (cast_arr, cast_count)),
Pre => cast_count in 0 .. Max_Votes-1,
and then not already_cast(paper)
      ...
Post => already_cast(paper)
and then cast_count = cast_count' old + 1);

```

```

procedure cast(paper : in barcode) is
begin
  cast_arr(cast_count) := paper;
  cast_count := cast_count + 1;
end cast;

```

Transformation Issues

- What do we **prove** at SPARK level?
 - **Not** necessarily all system invariants need to be re-proved in SPARK (already proved in Event-B)
 - Need to prove Ada is a **correct implementation** of the Event-B model
 - Some invariants might be required (e.g., well definedness)

Conclusions

- The SBB Event-B model
 - Modelled different **security properties**: Availability, confidentiality & integrity
 - Showed how we applied a refinement-based approach to model security properties
- Manual **Transformation** of Event-B Models to SPARK
- Identification of Translation Patterns
 - Applied to SBB & Tokeneer

Future Work

- What **additional assertions** are needed at SPARK level (invariants)
- Automatic **Code Generation**
 - Define a SPARK EMF Metamodel using XSD schema generated by GNAT
 - Event-B EMF to SPARK EMF Transformation

Thank you
Questions?