

Verifying System-level Security of a Smart Ballot Box

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Outline

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Motivation

- Application of refinement-based formal modelling in building a Correct-by-Construction secure system.
- Refinement of the availability property of secure systems.
- Overall Aim of case study: show how the Smart Ballot Box can be correctly implemented on capability hardware according to the system-level security specification.

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Case Study: Smart Ballot Box

- Key Function of SBB:
 - Ensures only valid ballot papers are cast in ballot boxes for later tabulation.

- Security Properties:
 - Confidentiality, integrity and availability.

Galois and Free & Fair. The BESSPIN Voting System (2019).

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Rigid Events and Parameters

- Event availability in Event-B
 - Determined by its enabledness condition.
 - Guard strengthening can affect event availability during refinement.
- Extend the notion of event **enabledness** to include parameters, given event *e* we define enabledness:

event e any p, q where G(p,q) then ... end

Enabled_p(e)
$$\stackrel{\text{\tiny def}}{=} \exists q . G(p, q)$$

• We call events we are interested in their **availability** with respect to *p* **rigid** events & *p* are the **rigid** parameters.

Rigid Events and Parameters (2)

Textual Representation: Event *e* must be enabled for any parameter *rp* satisfying *Enabled*_{rp}(*e*)

event [e] any [rp] op where Ga(rp, op) then...end

- Syntactic Rules:
 - 1. Rigid events can only be refined by *rigid* events
 - 2. The abstract *rigid* parameters must be retained in the concrete events
- In general, more rigid parameters can be introduced in later refinements, but they will only be relevant to *proof* in further refinements.

Preserving Availability through Refinement

- Preserve availability property through refinement by:
 - Proving that the concrete event does not strengthen the enabledness of the abstract event, we propose *enabledness* PO: *ENBL*

$$I(v), J(v, w), Ga(rp, oap, v) \vdash \exists ocp . Gc(rp, ocp, v, w)$$

event [ae]
any [rp] oap
where Ga(rp, oap, v)
then
// abstract actions
end

event [ce]
any [rp] ocp
where Gc(rp, ocp, v, w)
then
// concrete actions
end



Preserving Availability (2)

• In Event-B an abstract event can be refined by a group of concrete events ce_i ($i \in 1..n$).

ENBL PO can be generalized as follows where ocp_i and Gc_i are the concrete events and guards of ce_i

 \forall rp, oap . Ga(rp, oap) \Rightarrow \forall _i (\exists ocp_i . Gc_i(rp, ocp_i))

SBB System Model: Refinement Strategy

- o. Abstract level: Model an ideal voting system.
- 1. Model possible attackers' behavior by distinguishing between different types of ballot papers.
- 2. Introduce time and invalidate ballots with expired timestamps.
 - Time can be the subject of more attacks.
- 3. Data refine the voter information by encrypting ballots.
- 4. Ensure the legitimacy of ballots through the Message Authentication Code (MAC).



SBB Model: Abstract Level

- Events: create_ballot, cast ballot, invalidate_ballot
- Model an ideal voting system
 - Each voter can have at most one legitimate ballot

ballots \in VOTER \Rightarrow VOTE

– The cast ballots must be legitimate

 $cast \subseteq ballots$



First Refinement: Ballot types



- Possible attacks
 - Attacker create ballot/duplicate valid ballot ..
- Model the main security properties of SBB
 - 1. Accept all valid ballots
 - 2. Reject invalid ballots

First Refinement: Availability Property

- Availability property: Ensure valid ballots are not blocked from being cast.
- Availability property is captured by the guard of the relevant events.
 - Specify *cast_paper* as **rigid event** with *paper* as
 - rigid parameter.

event [cast_paper] refines cast_ballot
any [paper] where
@valid-paper: paper ∈ valid_papers
then
// actions for casting a ballot
end

Second Refinement: Time & Availability Enc

Encoding **ENBL** PO as a theorem

Southampto

event cast_paper refines cast_paper any paper where

@typeof-paper: paper ∈ papers

// paper not already expired

// copy not already cast
// copy not already spoiled

// paper is not illegitimate

then

// cast the paper actions
end

theorem @accept-valid-paper: \forall paper \cdot paper \in valid_papers \Rightarrow // paper not already expired paper time(paper) ≥ current time – expiry duration // copy not already cast ∧ paper voter(paper) ∉ paper_voter[cast_papers] // copy not already spoiled ∧ (\forall sp · sp ∈ spoiled_papers ⇒ paper_voter(paper) # paper_voter(sp) V paper_vote(paper) ≠ paper_vote(sp) V paper_time(paper) ≠ paper_time(sp)

// paper is not illegitimate∧ paper ∉ illegitimate papers

Third Refinement: Ballot Encryption

- Introduce encryption to prevent SBB from accessing the voter's information.
 - Apply data refinement to replace *paper_voter* and *paper_vote* with encrypted ballot
- Prove *ENBL* PO due to *cast_paper* guards update as a result of refinement.

```
theorem @accept-valid-paper:
\forall paper \cdot paper \in valid_papers \Rightarrow
paper_time(paper) ≥ current_time -
expiry duration
// copy not already cast
∧ paper_encrypted_ballot(paper) ∉
paper_encrypted_ballot[cast_papers]
// copy not already spoiled
\land (\forallsp \cdot sp \in spoiled papers \Rightarrow
paper_encrypted_ballot(paper) #
paper_encrypted_ballot(sp) V
paper_time(paper) # paper_time(sp)
∧ paper ∉ illegitimate_papers
```

Fourth Refinement: Ballot Authentication

- Introduce MAC to check the legitimacy of the source issuing the ballot.
 - We assume the attacker does not know the secret key; therefore, it is crucial to ensure the secrecy of this key.

```
@mac-legitimate_papers: ∀paper · paper ∈ legitimate_papers ⇒
paper_mac(paper) = MACAlgorithm(
paper_time(paper) ↦ paper_encrypted_ballot(paper) ↦ MACKey
)
```

• The guards of *cast_paper* and The *ENBL* PO will be updated accordingly.

Conclusions and Future Work

- Availability property of an event can be ensured through refinement by preserving the enabledness of its corresponding refined events.
 - A general PO (*ENBL*) that can be applied to any event with rigid parameters is provided.
- Future Work
 - Semantics model to justify the soundness of the rigid property of events
 - Tool support for the ENBL PO in Rodin: CamilleX



Thank you

Questions?

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Visit <u>https://hd-sec.github.io</u> for more information on the HD-Sec project.